

# ME/AER 676 Robot Modeling & Control

## Spring 2023

### Planning As (Graph) Search

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# Graphs

A Graph  $G = (V, E)$  Consists of two lists (sets):

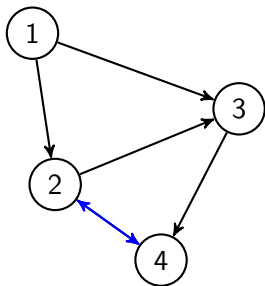
- ▶ Vertices  $V$

On the right,  $V = \{1, 2, 3, 4\}$

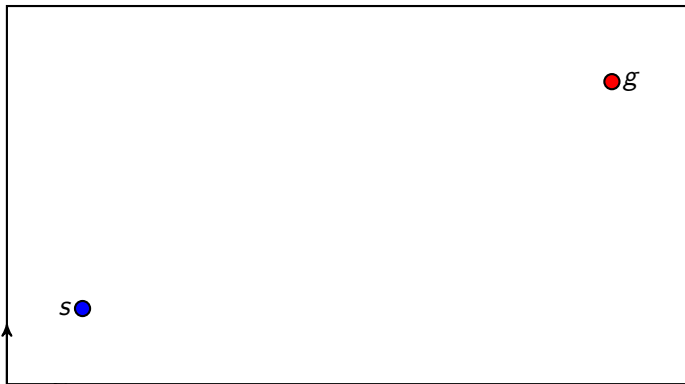
- ▶ Edges  $E \subseteq V \times V$

On the right,  $E =$

$\{(1, 2), (1, 3), (2, 3), (2, 4), (3, 4), (4, 2)\}$

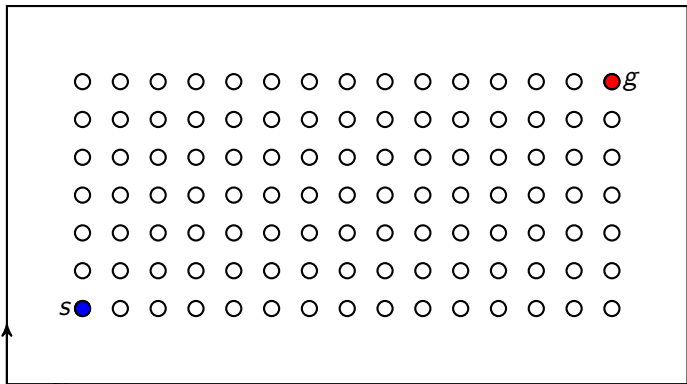


# Graph Search



$s = \text{start}$ ,  $g = \text{goal}$

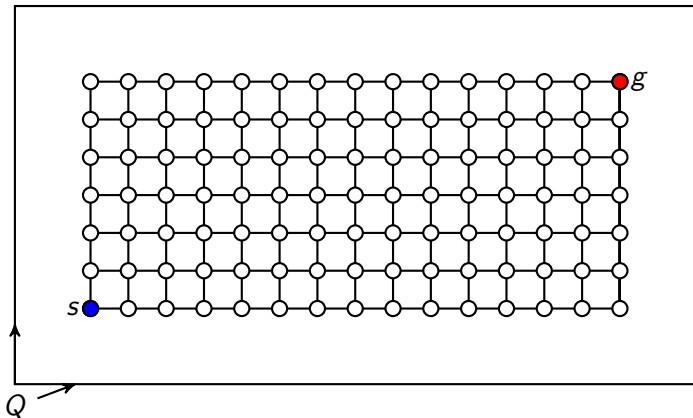
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Define nodes

# Graph Search



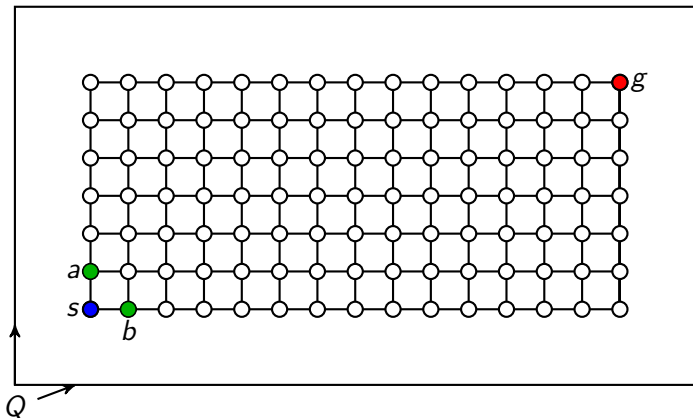
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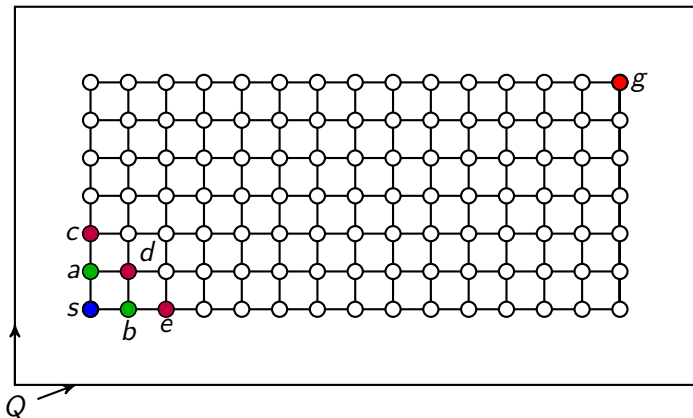
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Define nodes

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Let's focus on the nodes close to  $s$

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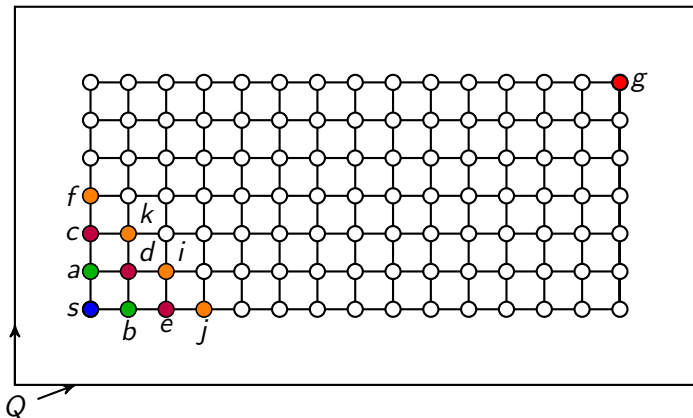
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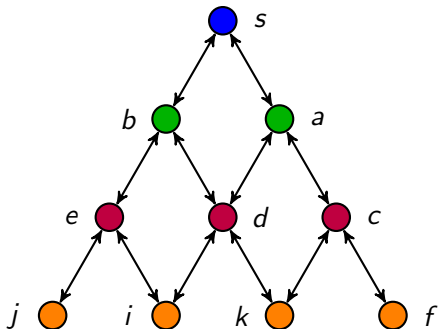
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## Graph As Seen From Start



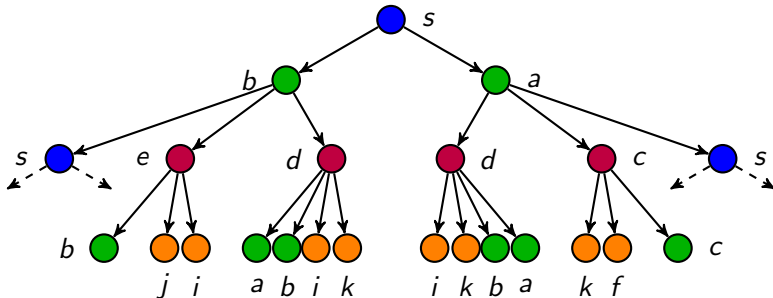
The undirected edges are equivalent to directed edges going forwards and back between two nodes.

We will use a search tree to **traverse** paths in this undirected graph.

# Search Tree

The root of the search tree is node  $s$ .

This tree depicts the possible paths starting from  $s$ .



Each 'layer' contains the nodes reached after taking a step from a node in a preceding layer.

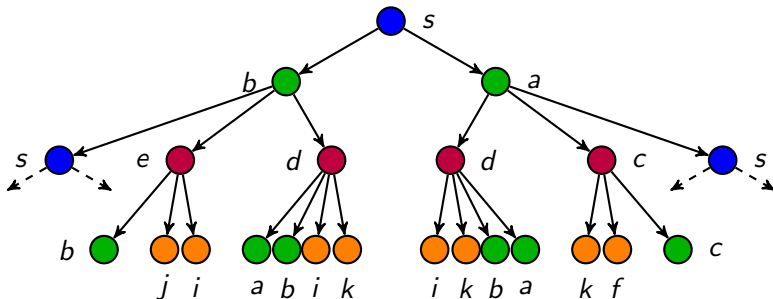
Each layer has a depth, starting at 0 for  $s$  on top.

Node  $s$  appears at depth 0 and 2;  $b$  at 1 and 3, etc.

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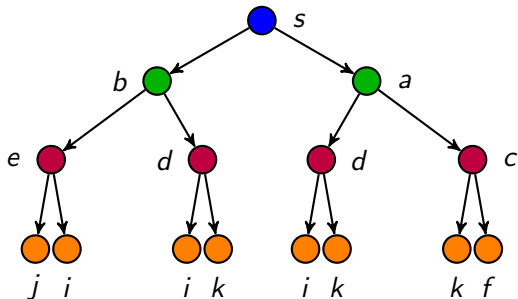
The actual tree is much larger, with infinite depth.

For example, the tree recursively repeats itself whenever  $s$  appears at some depth.

# Search Tree

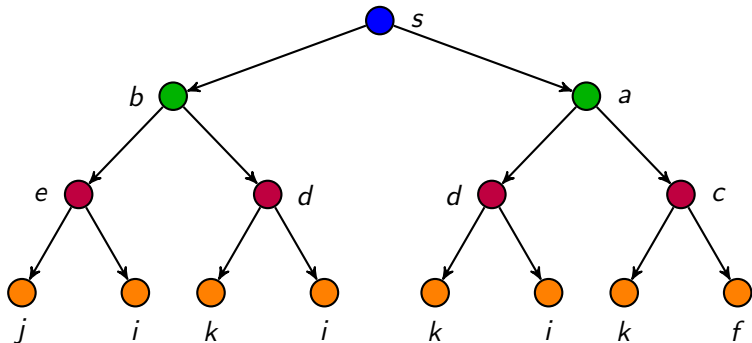
The root of the search tree is node  $s$ .

This tree depicts the possible paths starting from  $s$ .



If the edges in the graph were directed, we would never return to nodes located at smaller depths, simplifying the tree.

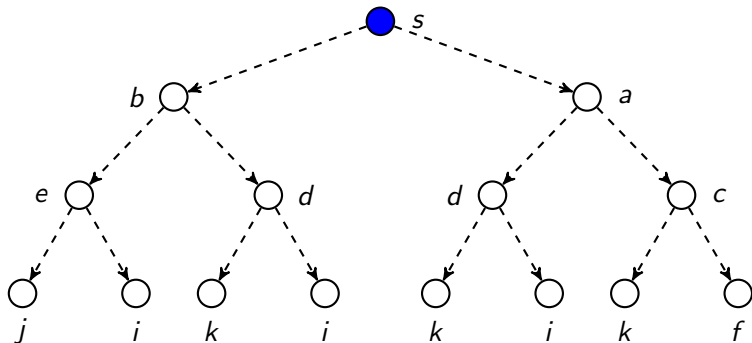
# Search Tree



If a path exists from  $s$  to  $g$ , it is one of the possible paths in this search tree.

The key question is: how do we enumerate all the possible paths starting from  $s$ ?

# Search Algorithms



Search algorithms avoid representing all paths, or the entire tree, at one go.

Instead, the search tree is incrementally traversed.

Start with root node, and no other nodes/edges.

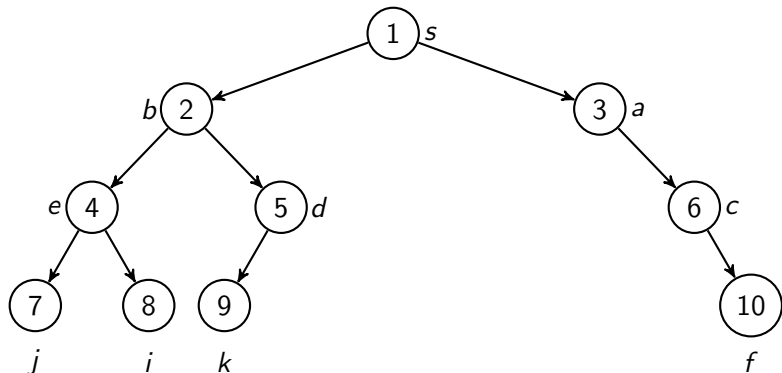
# Search Algorithms

Search algorithms differ in the order in which paths are evaluated, which boils down to the order in which unvisited successor states are checked for being the goal.

- ▶ Uninformed Search: Only know the successors at each node
  - ▶ Breadth First Search: Adds successors to end of list (Queue: first in first out)
  - ▶ Depth First Search: Adds successors to start of list (Stack: first in last out)
- ▶ Informed Search: Use an estimate of potential value of a node to order unvisited nodes (neither FIFO nor FILO).



# Breadth-First Search



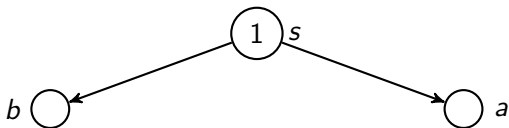
Order of visiting nodes under BFS. 'Row-by-row'

# Breadth-First Search

①<sub>s</sub>

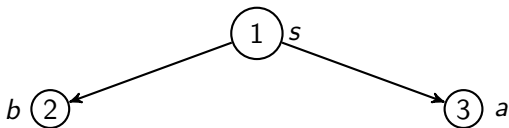
What the search tree looks like as it is built

# Breadth-First Search



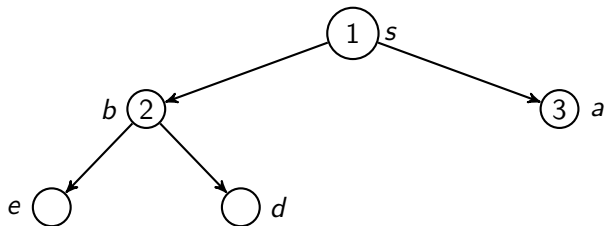
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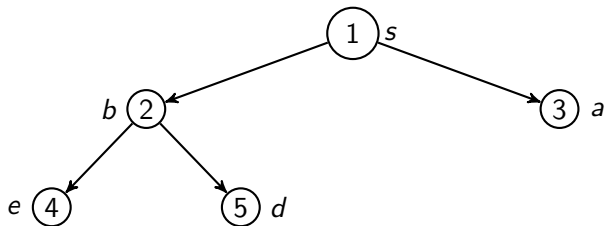
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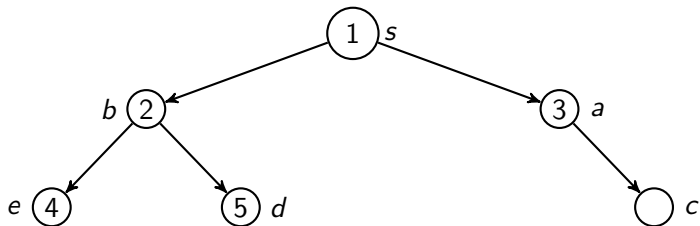
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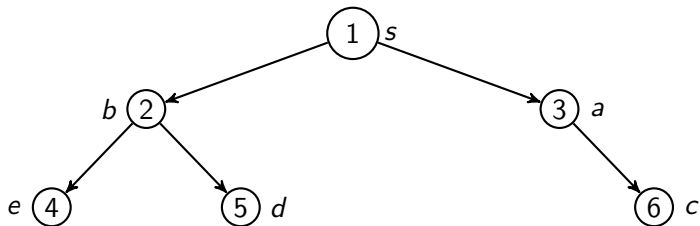
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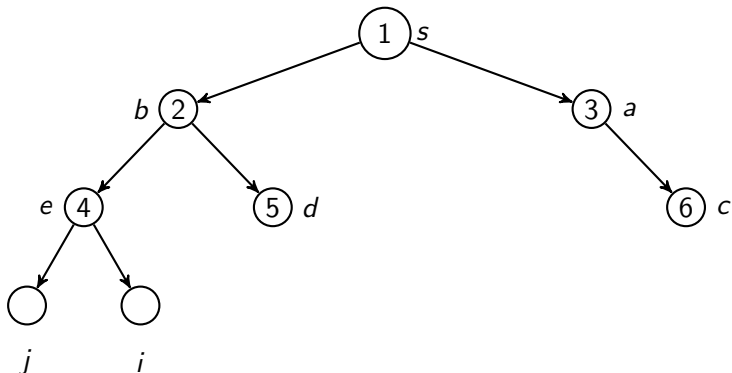
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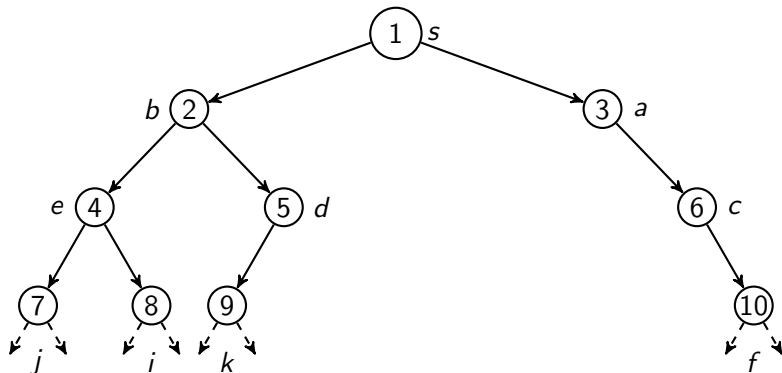


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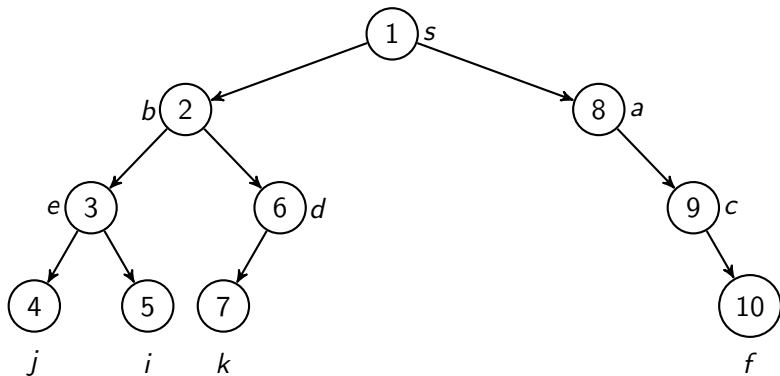
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# Depth-First Search



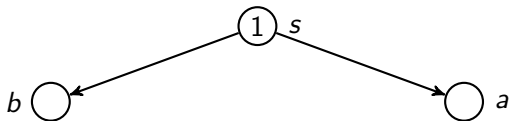
Order of visiting nodes under DFS. 'Column-by-column'

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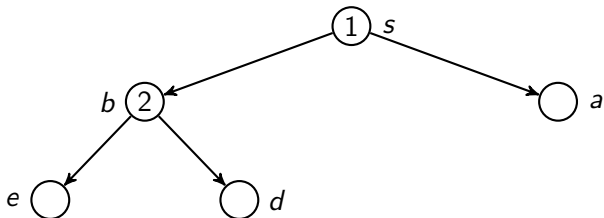
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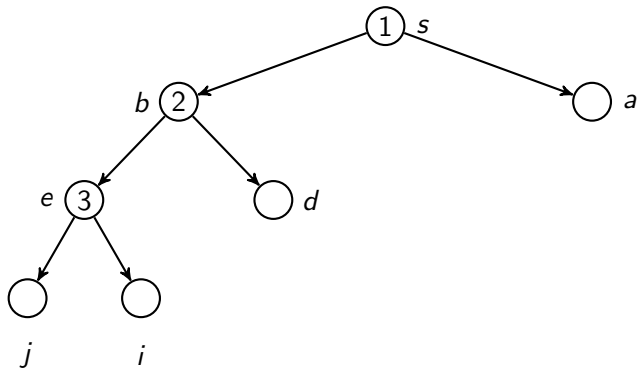
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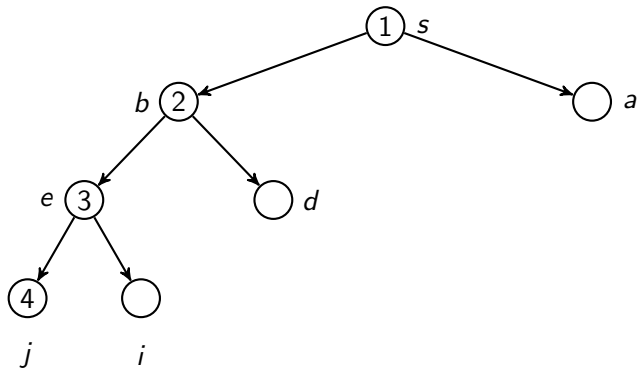
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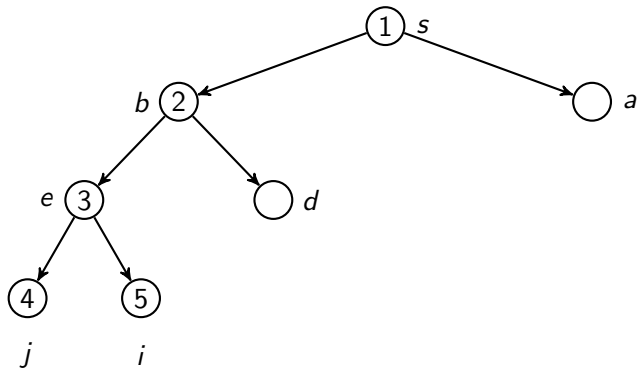
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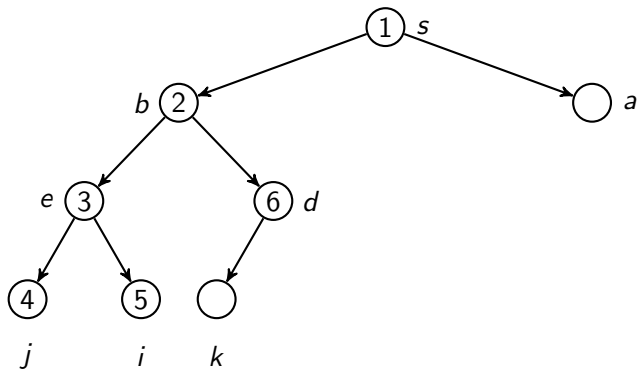


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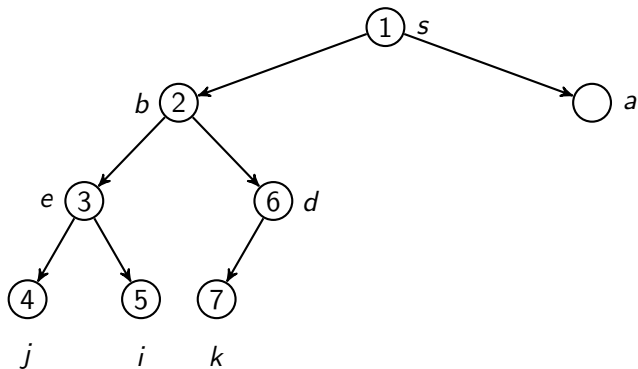
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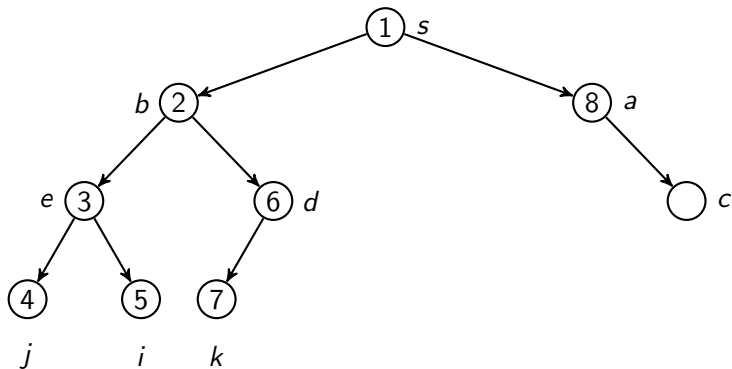
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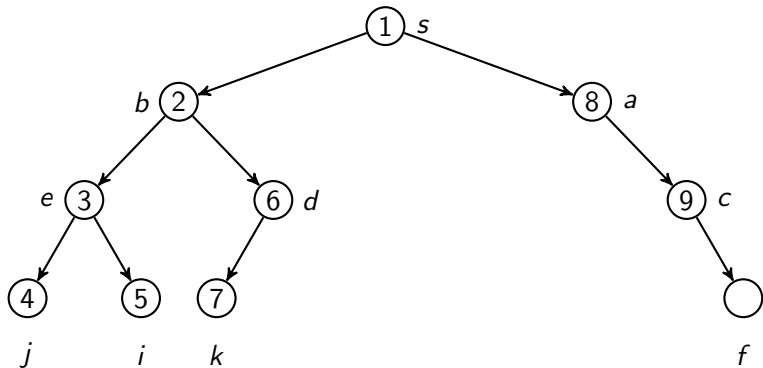
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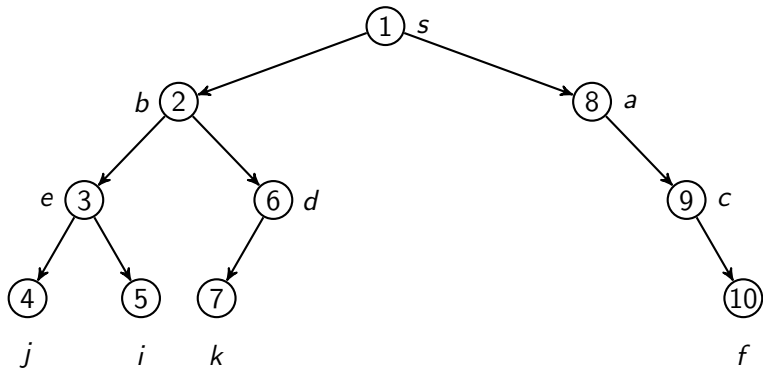
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- ▶ Some estimate of how likely a node will lead to the goal would be a useful way to break ties
- ▶ This idea leads to informed search, where the estimate for each node is provided by a *heuristic* function

## Three IS Algorithms

There are two quantities we can assign to a node:

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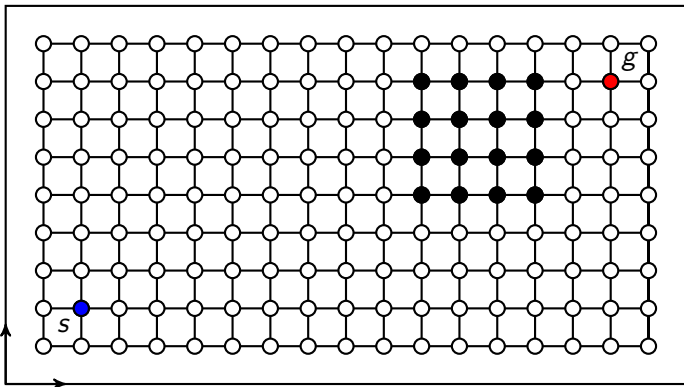
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1. Best-First Search: order using  $h(n)$
2. Dijkstra's algorithm: order using  $c(n)$
3. A\*: order using  $c(n) + h(n)$ . Guaranteed performance when

$$h(n) \leq h^*(n).$$

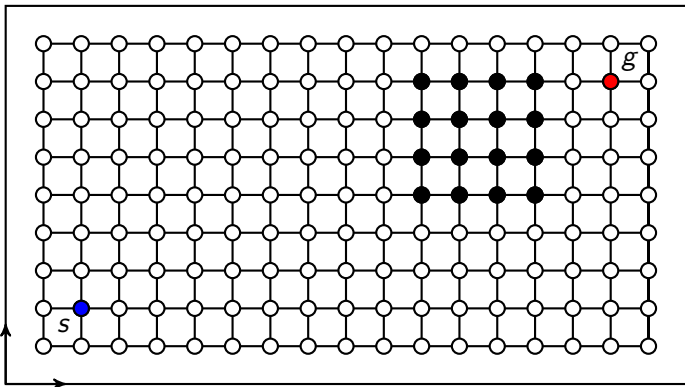


# Motion Planning Discrete Space



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While we 'see' the graph all at once, our algorithms deal with this graph using a search tree